

Unified Forwarding with Segment Routing

Mohan Nanduri
August 3, 2014



Some Challenges...

- Control plane scale and complexity
- Many protocol and features, also many bugs
- Forwarding table size and FIB capacity
- Programmatic control over BGP policy
- Hyper scale datacenters
 - Many encapsulations
 - Lot of links
 - Commodity hardware
 - IP address mobility
- Inter-DC vs External workloads
 - Latency-sensitive vs bulk
 - Scheduled vs unscheduled

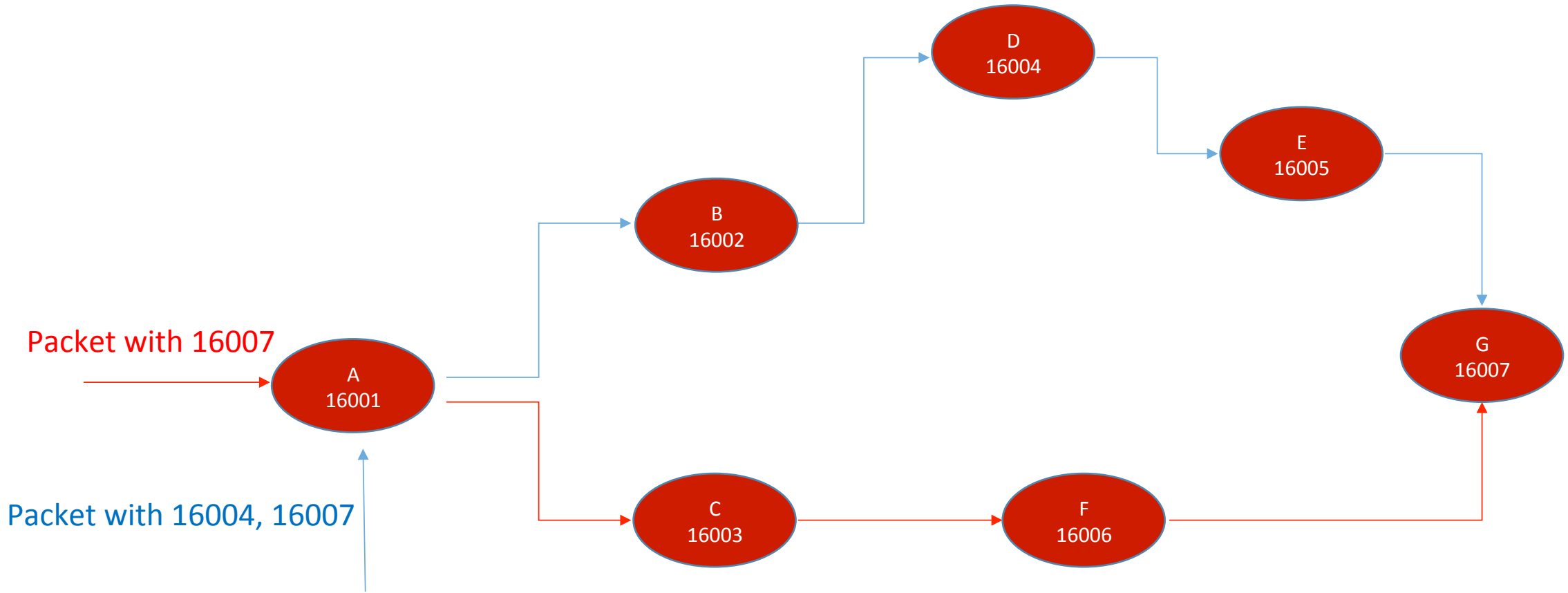
What's the goal?

- Reduce complexity and state in the network
 - Number of protocols
 - Control plane state
- Unified forwarding plane for core, edge and datacenter
- Programmatic interface to the network
- Reduce feature dependencies
- Leverage commodity hardware across network layers
- Reduce FIB table size

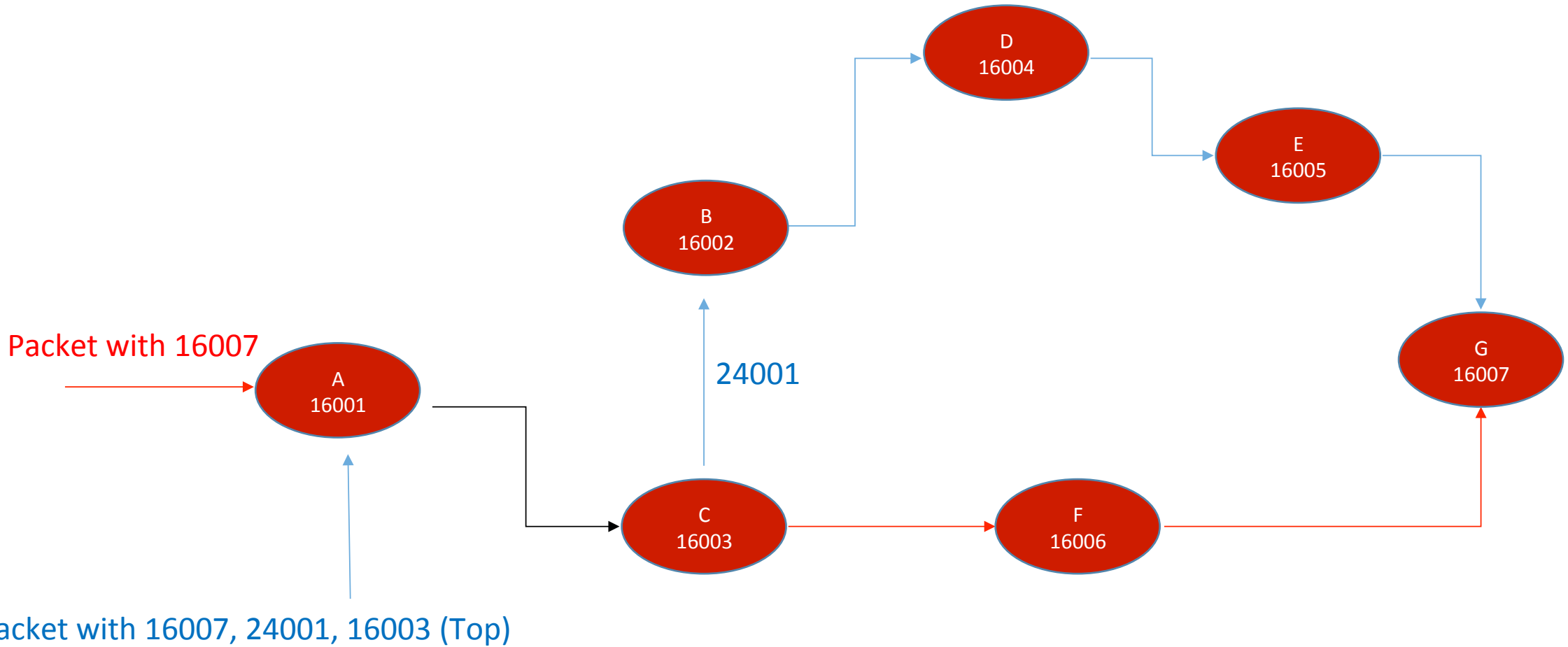
Segment Routing Primer

- Label-based source routing aka SPRING
- Uses existing MPLS data plane
- IGP floods labels throughout the SR domain
- Node-SIDs – Devices are configured with globally significant labels
 - Installed on all devices in the domain
- Adjacency-SIDs - Each SR router generates locally significant link labels
 - Installed only locally
 - They are still flooded to the entire domain

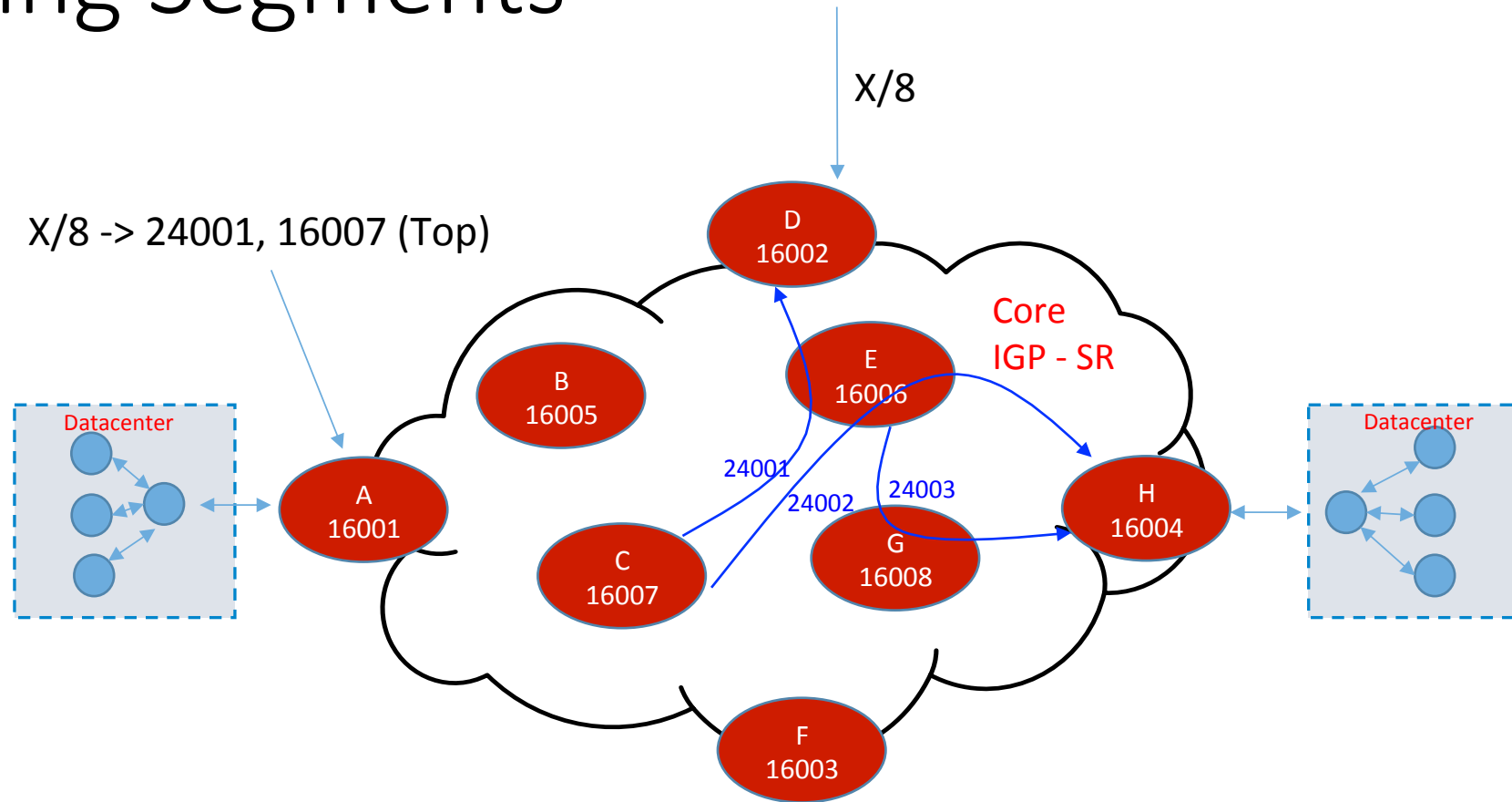
Use of Node SIDs, label distribution via IGP



Use of Adjacency SIDs and Shortest Path Override

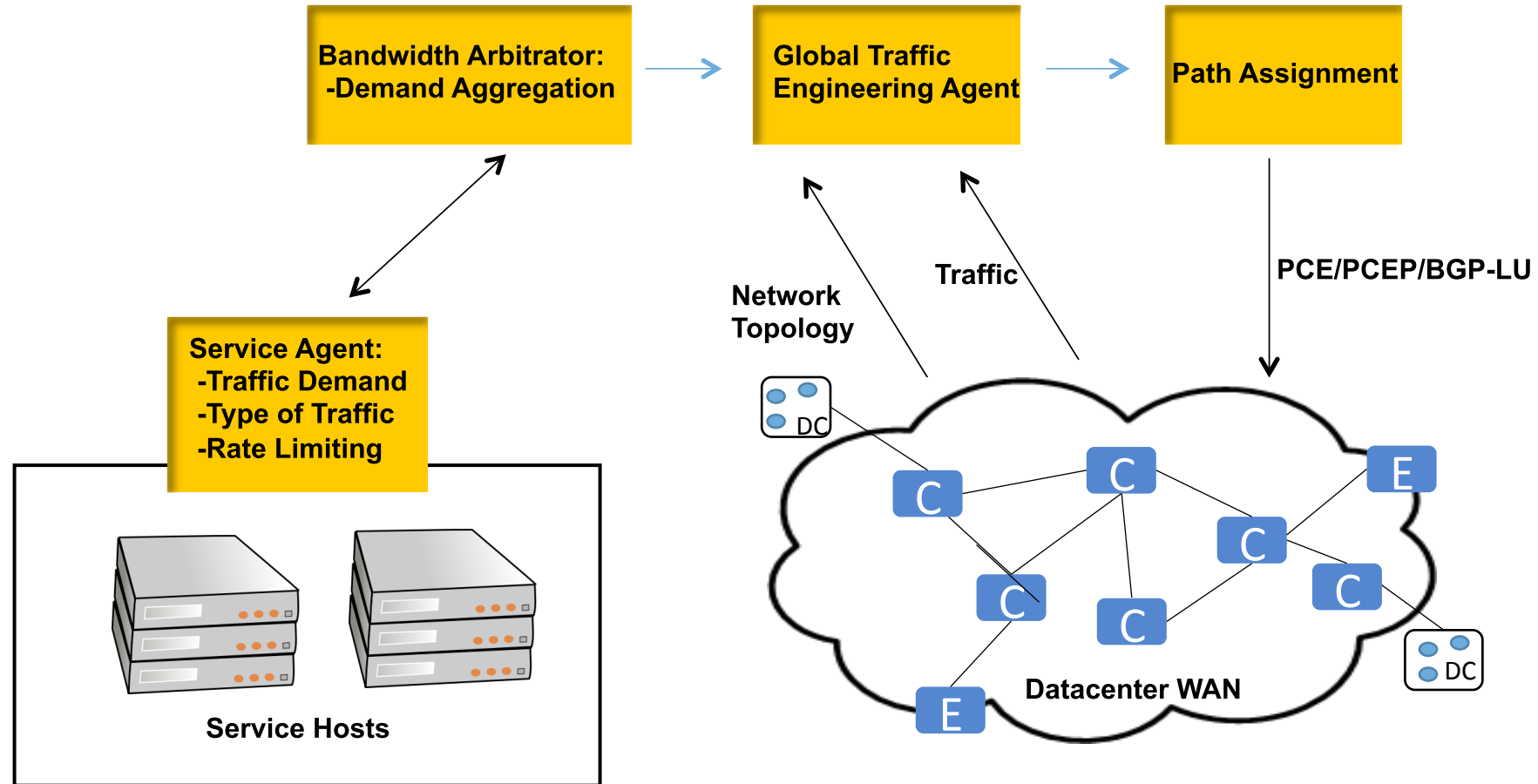


Binding Segments

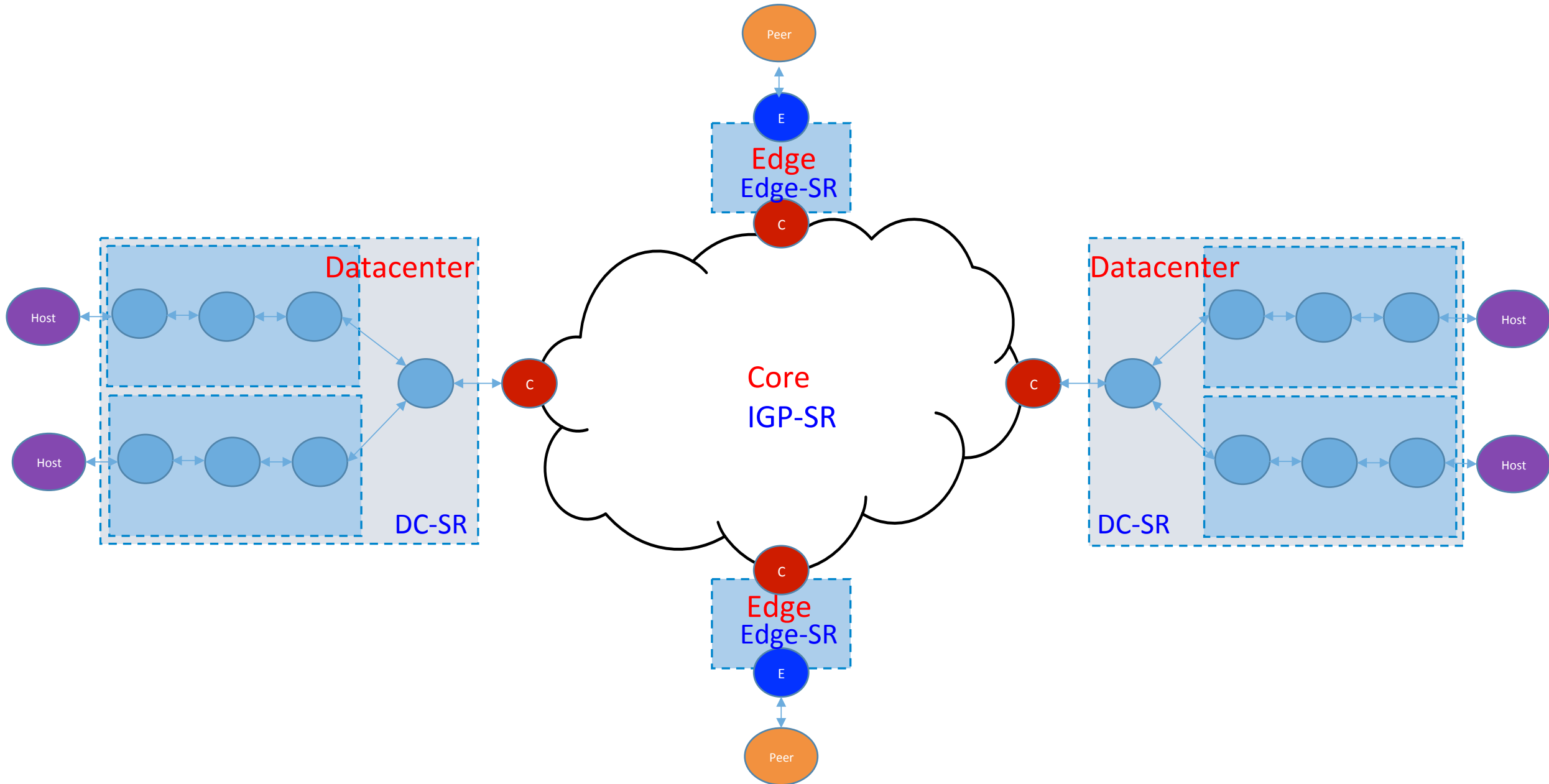


- Binding SID represents a kind of a tunnel, pop Binding SID, push one or more labels

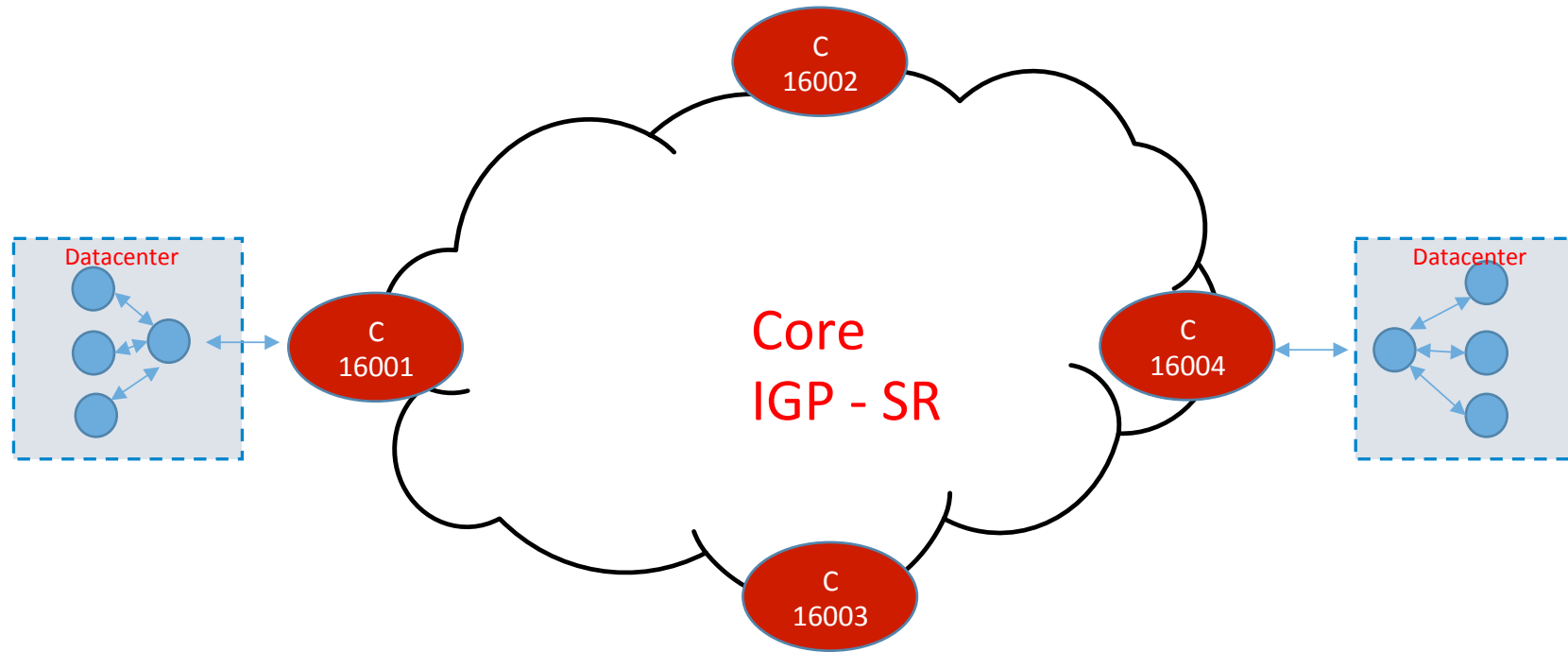
Software Driven Network



Unified Forwarding – Core, Edge and Datacenter



Unified Forwarding – Core



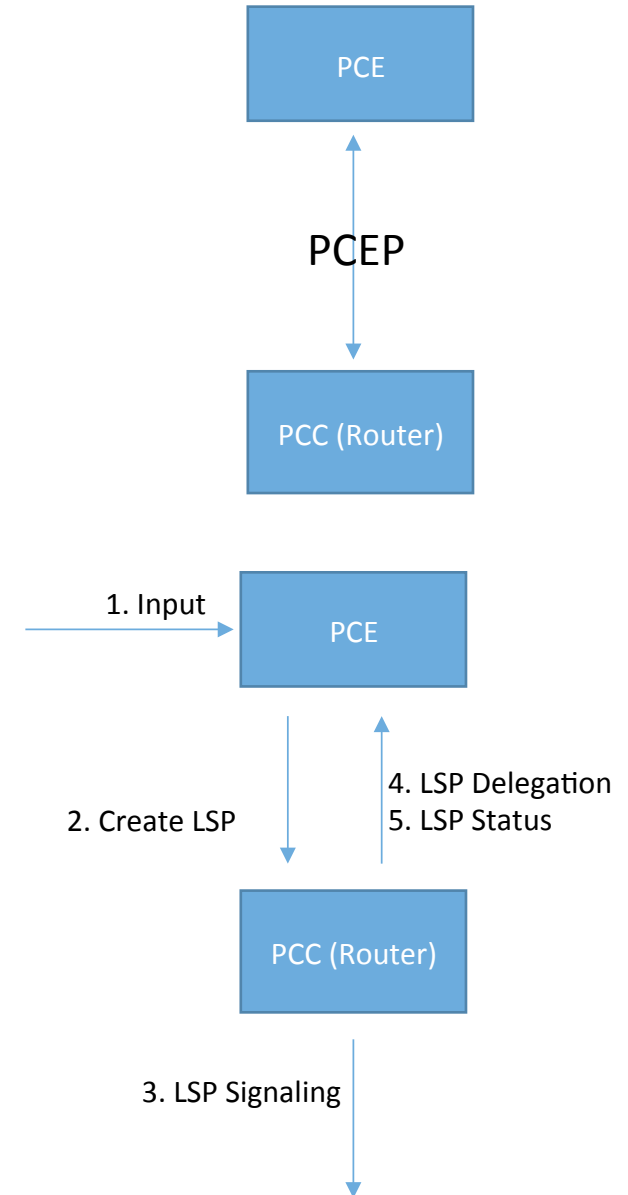
- SR Extensions for IGP
- PCE/PCEP with SR-TE
- BGP-LU

Core Traffic Engineering

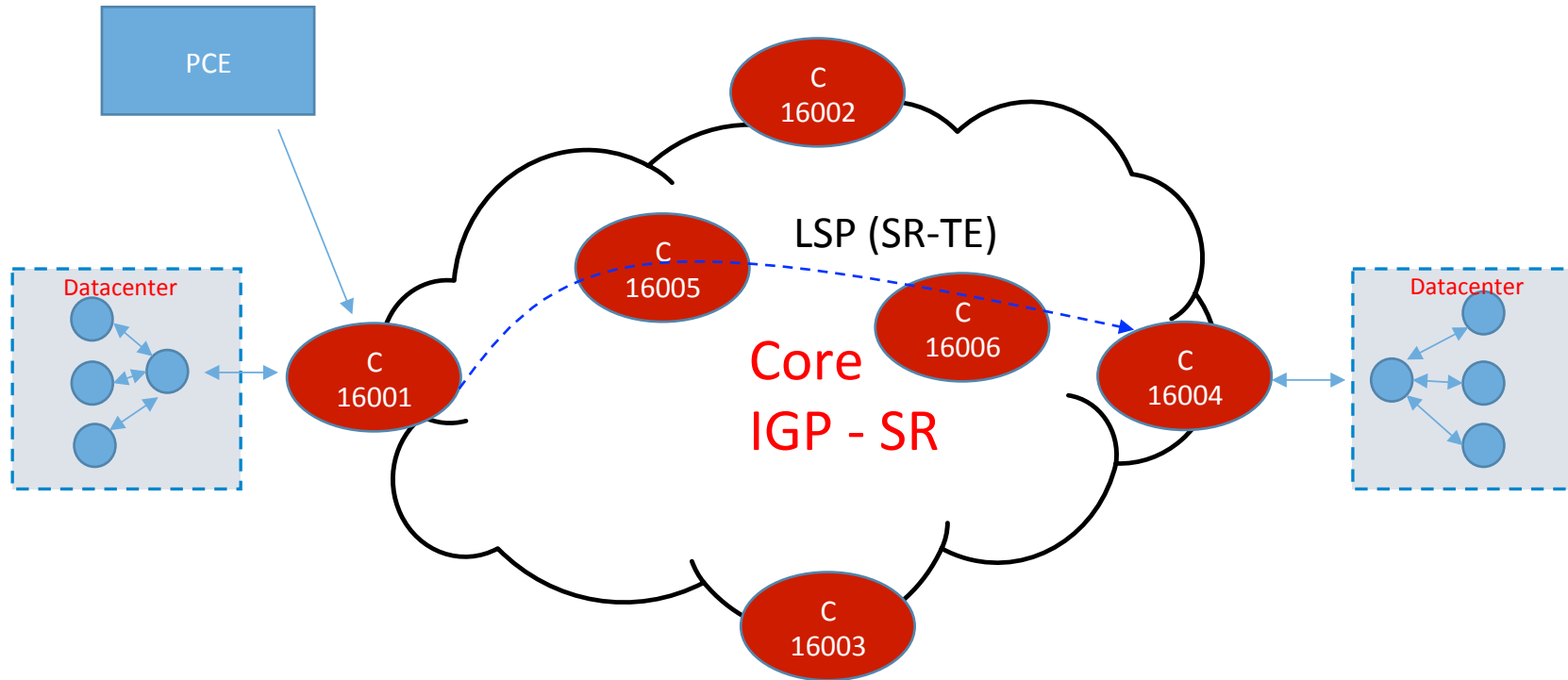
- PCE/PCEP
 - SR-TE
- BGP-TE
 - BGP-LU

PCE/PCEP with SR-TE

- Stateful PCE components
 - Path computation element (PCE)
 - Path computation client (PCC)
 - Path computation element communication protocol (PCEP)
- Offline computation and program path using PCEP



PCE/PCEP with SR-TE



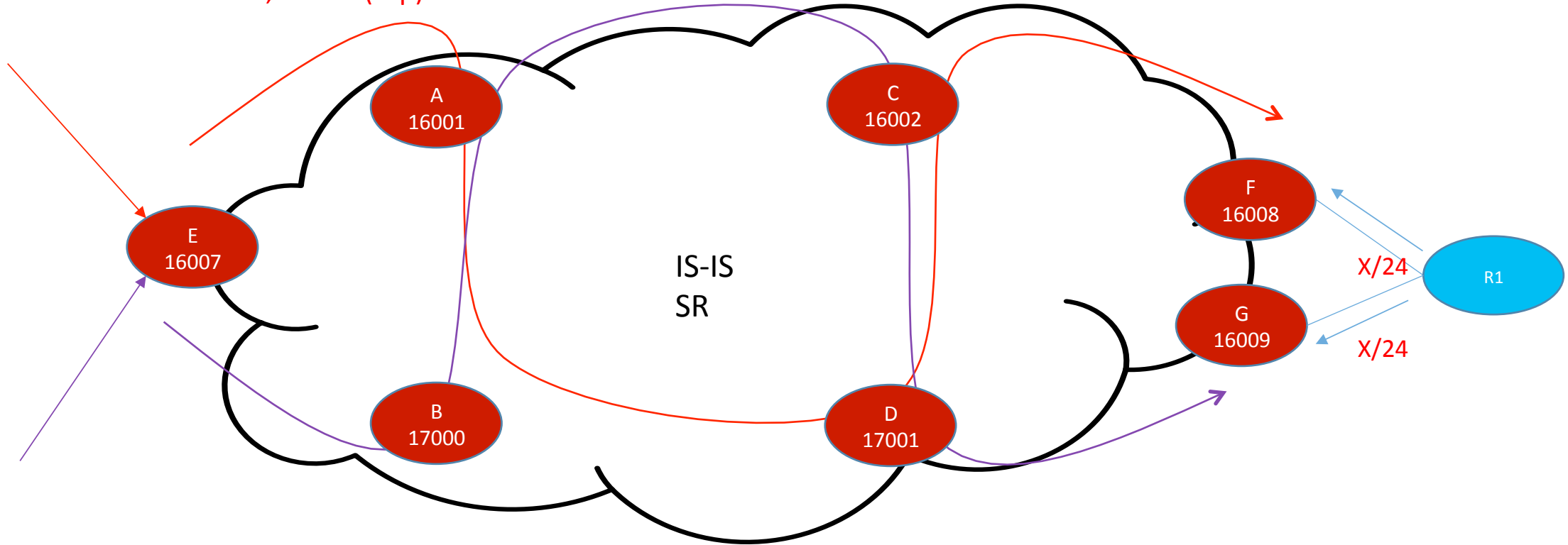
- Offline computation and PCE/Controller will program LSPs

BGP Traffic Engineering (BGP-TE)

- X/8 → nhop address, label stack 1, Link Bandwidth 1
nhop address, label stack 2, Link Bandwidth 2
- With Binding SID, any label in the stack can be a binding segment label

BGP-TE Load Balancing

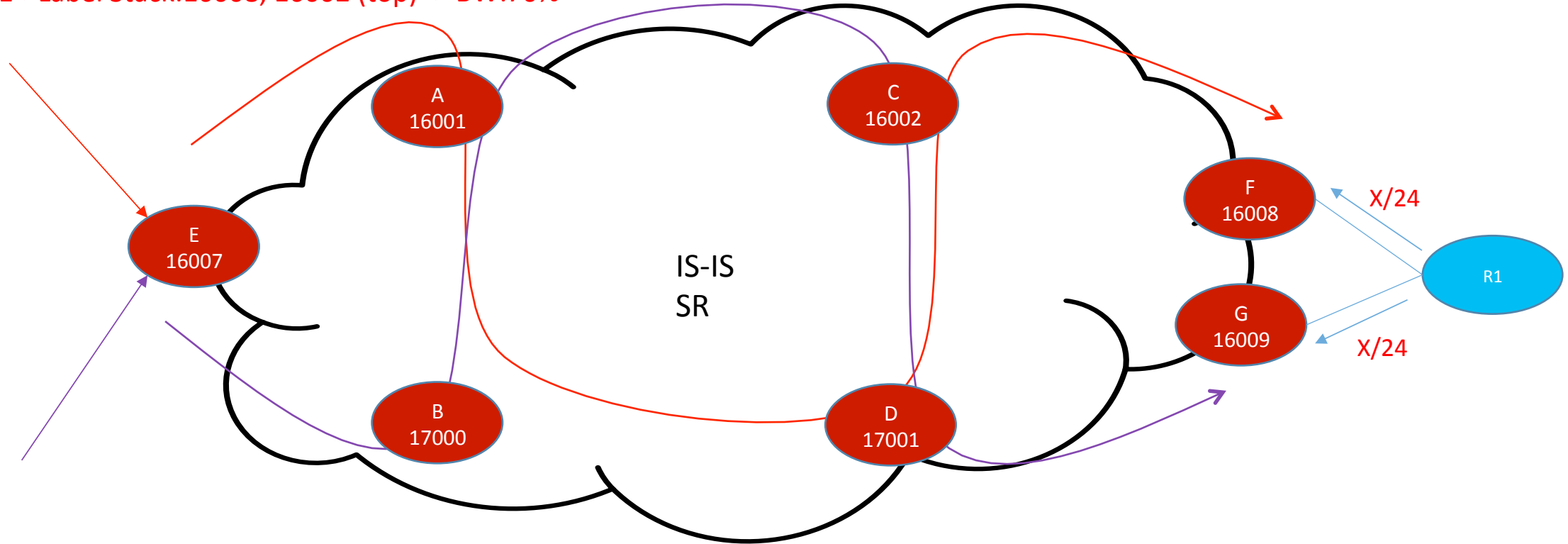
X/24 -> nh:if1->Label Stack:16008, 16002 (top) -> BW:50%



X/24 ->nh:if2->Label Stack:16009,17001 (top) -> BW:50%

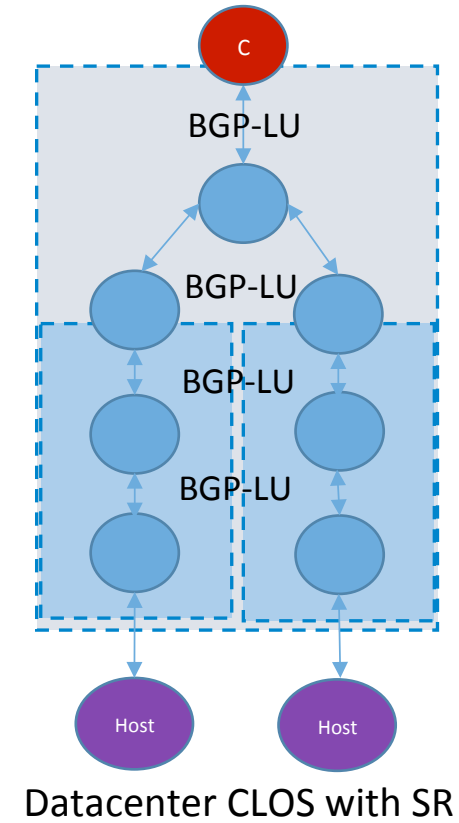
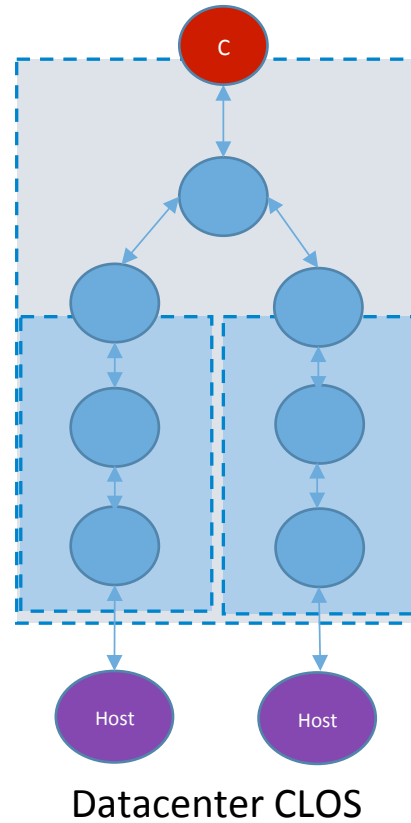
BGP-TE Unequal Cost Load Balancing

X/24 -> nh:if1->Label Stack:16008, 16002 (top) -> BW:70%



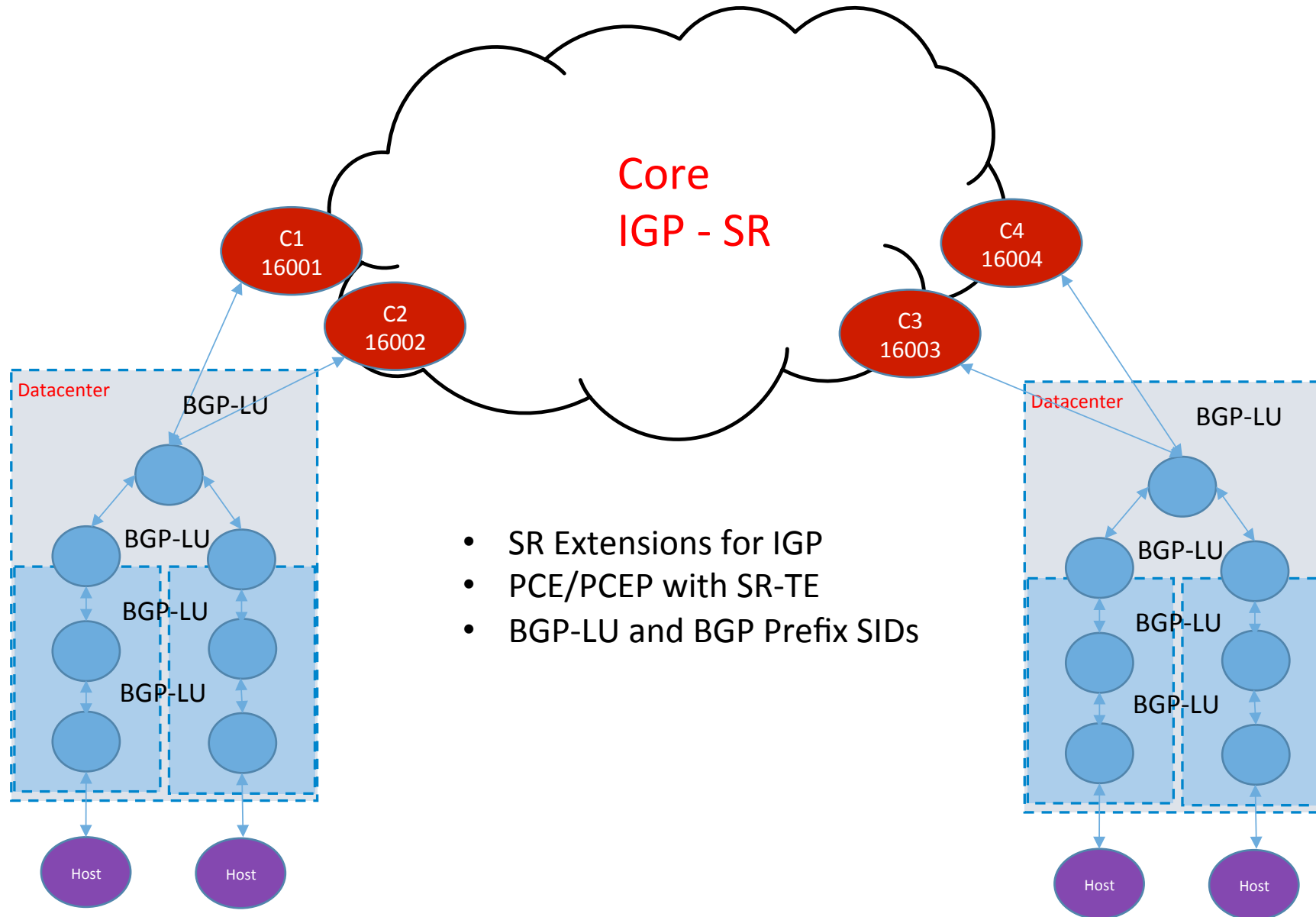
X/24 ->nh:if2->Label Stack:16009,17001 (top) -> BW:30%

Unified Forwarding – Datacenter

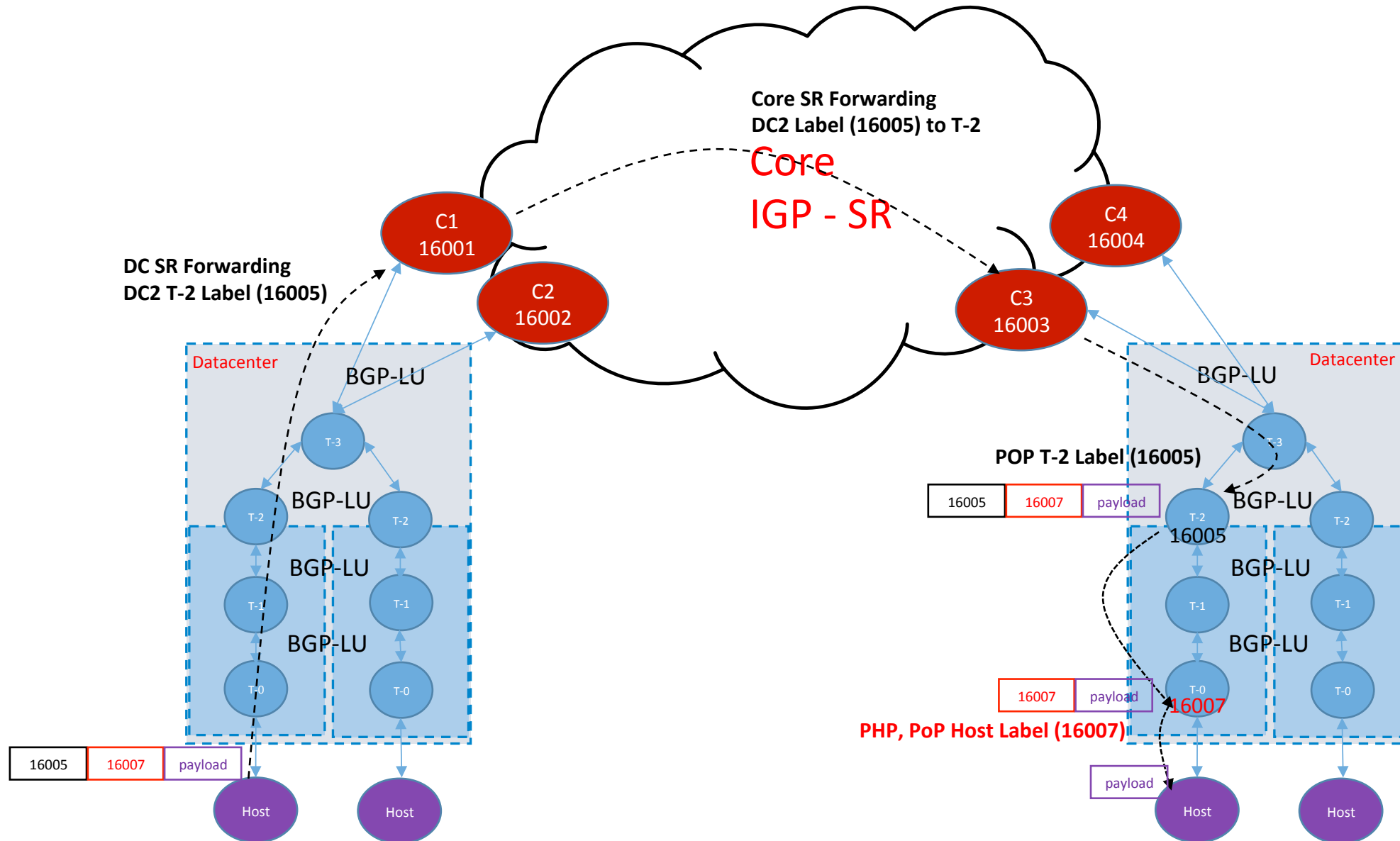


- BGP Prefix-SIDs
- BGP-LU

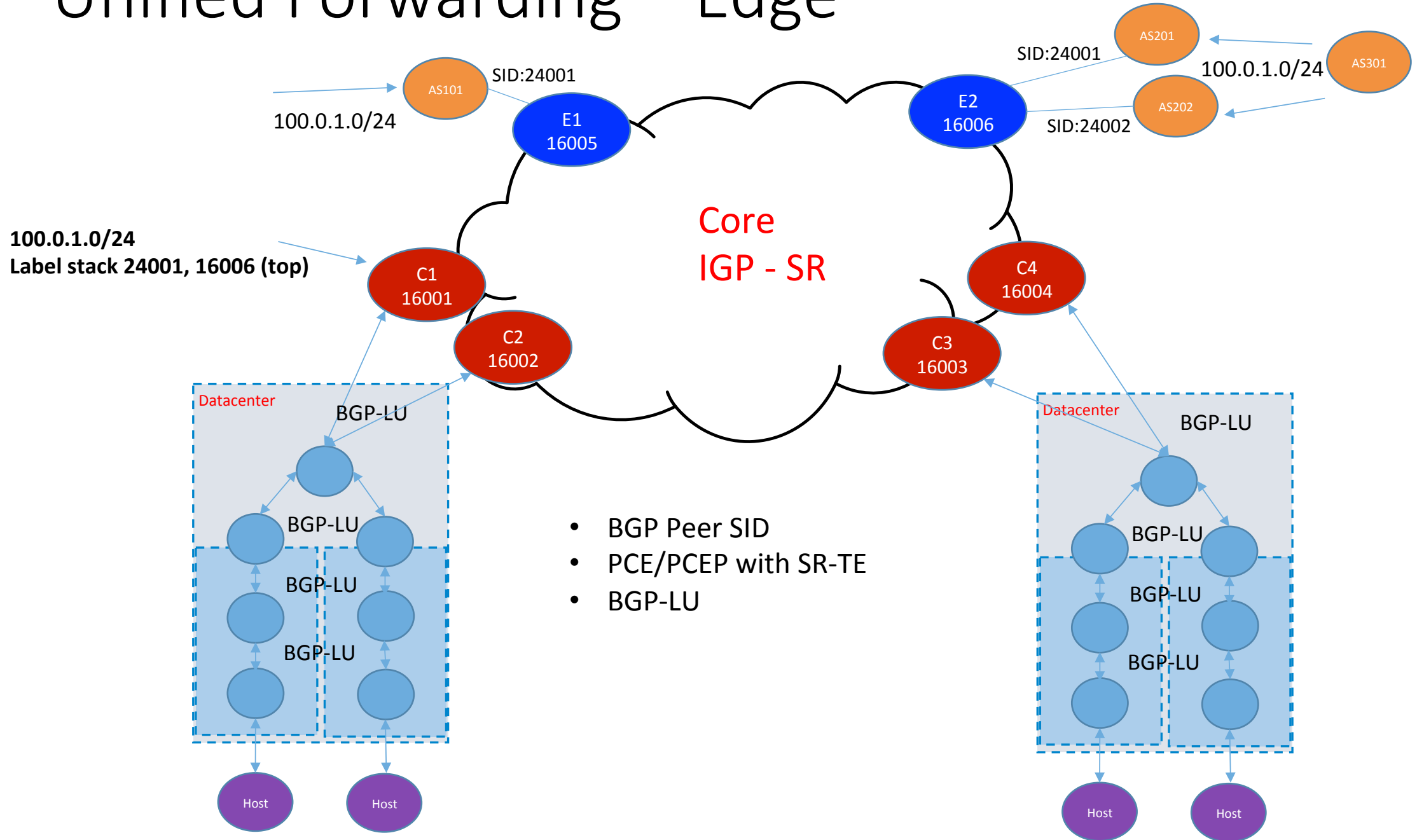
Unified Forwarding – Core and Datacenter



Unified Forwarding – Core and Datacenter



Unified Forwarding – Edge

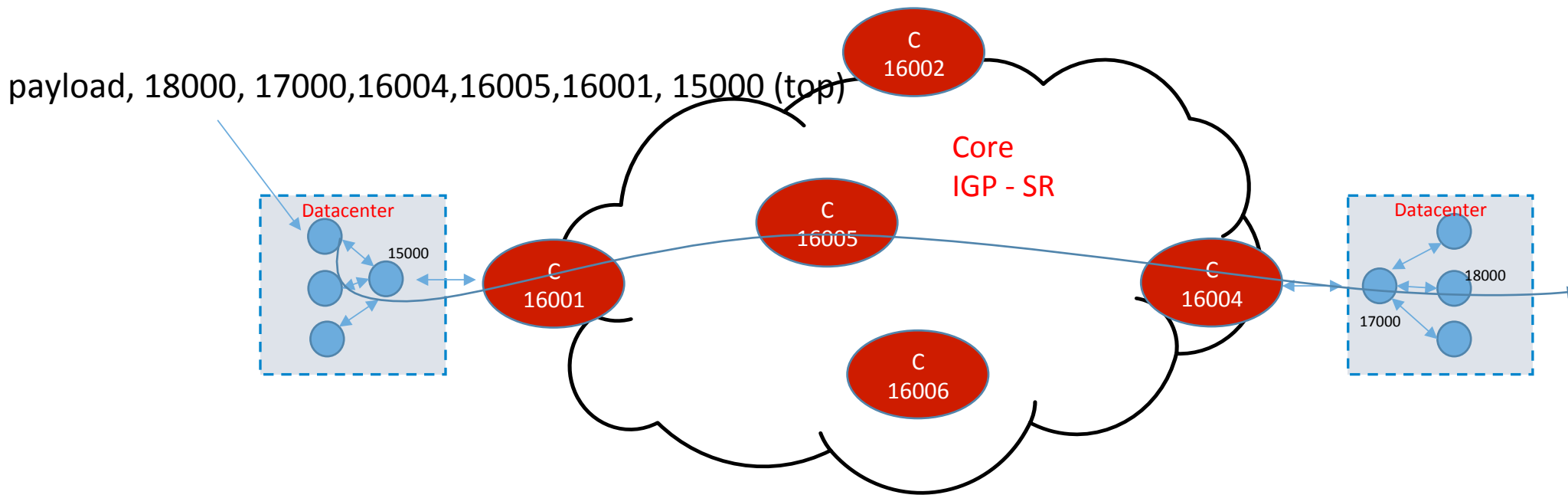


Unified Forwarding – Edge

- Granular traffic engineering
 - Prefix based
 - AS based
 - Overriding egress next-hops
 - Performance based routing, diverting traffic based on performance and load
- Centralized control
 - peering egress control
 - Policy engine, reduces custom configuration and standardize peering policies
- Security
 - Remote black hole triggering
 - Inject flowspec rules

Data plane Monitoring

- Simplifies end-to-end monitoring
- Construct and signal probe packets for data plane health check
- Construct paths without creating state in the network



Useful URLs

- <http://www.segment-routing.net/>
- <http://datatracker.ietf.org/wg/spring/documents/>
- <http://datatracker.ietf.org/wg/idr/documents/>
- <http://datatracker.ietf.org/wg/pce/documents/>

